

Ecole polytechnique fédérale de Zurich Politecnico federale di Zurigo Swiss Federal Institute of Technoloav Zurich

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# Algorithms, Probability, and Computing Fall 08 Mid-Term Exam

#### Candidate:

First name:	
Last name:	
Registration Number:	

I attest with my signature that I could have taken the exam under regular conditions and that I have read and understood the general remarks below.

Signature: .....

#### General remarks and instructions:

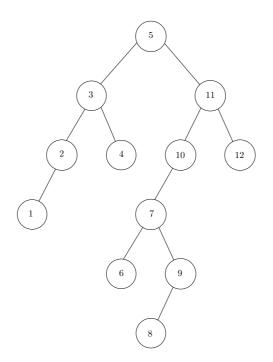
- 1. You can solve the 5 exercises in any order. You should not be worried if you cannot solve all the exercises (best grade with at least 48 points, i.e. 4 exercises).
- 2. Check your exam documents for completeness (1 cover sheet and 3 sheets containing 5 exercises).
- 3. Immediately inform an assistant in case you are not able to take the exam under regular conditions. Later complaints are not accepted.
- 4. Pencils are not allowed. Pencil-written solutions will not be reviewed.
- 5. No additives allowed.
- 6. Attempts to defraud yield to immediate exclusion from the exam and can have judicial consequences.
- 7. Provide only one solution to each exercise. Cancel invalid solutions clearly.
- 8. All solutions must be understandable and well-founded. Write down the important thoughts in articulative sentences and keywords. Unfounded or incomprehensible solutions will not be awarded. You can write your solution in English or German.
- 9. Make sure to write your name on all the sheets.

	achieved points (maximum)	reviewer's signature
1	(12)	
2	(12)	
3	(12)	
4	(12)	
5	(12)	
Σ	(60)	

In this exercise we consider random search trees for n keys.

- (a) What is the probability that the element with rank i is a common ancestor of the element with rank 1 and the element with rank 4? (You will have to distinguish between several cases)
- (b) What is the probability that the element with rank i is on the path from the root to the smallest element?

Let T be a search tree for  $\{1, 2, ..., n\}$ . The picture below shows an illustration

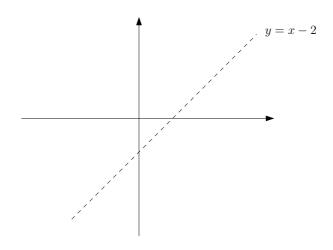


- (a) By a fancy triple we denote a triple  $\{i, i + 1, i + 2\}$  such that *i* is an ancestor of i + 2 and i + 2 is an ancestor of i + 1. (In the above picture we have two fancy triples, i.e.,  $\{5, 6, 7\}$  and  $\{7, 8, 9\}$ ) What is the expected number of fancy triples in a random search tree for  $\{1, 2, ..., n\}$ ?
- (b) Let  $k \le n$  be an integer. We say that a tree contains a k-tail if i is the child of i + 1 for every i with  $1 \le i \le k 1$ . E.g. the tree in the above illustration contains a 3-tail (and therefore also a 2-tail and a 1-tail). What is the probability that a random search tree for  $\{1, 2, ..., n\}$  contains a k-tail?

Consider the line l: y = 2. Let S be the set of non-vertical lines  $l' \neq l$  such that

- (i) l' intersects l and
- (ii) the x-coordinate of the intersection point is at least 1.

Draw the set of points whose dual is in S!

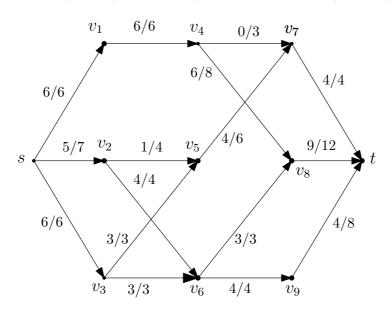


In this exercise we want to obtain a suitable preprocessing/query algorithm for efficiently deciding whether a point is inside a convex polygon, including giving a short certificate. Let P be a convex polygon with n vertices. We are interested in the following query. The input is a point q in the plane. The output is

 $\begin{cases} 3 \text{ vertices of } P \text{ such that the triangle spanned by these } 3 \text{ vertices contains } q, & \text{if } q \text{ is in } P \\ \text{no,} & \text{otherwise} \end{cases}$ 

Describe a data structure that stores P in O(n) space so that the query can be processed in time  $O(\log n)$ . **Hint:** If a point q is in P then lowest vertex  $p_{\perp}$  of P is part of a triangle containing q.

Consider the below network (G, s, t, c) with flow f (the edge labels are flow/capacity).



(a) What is the value of f?

 $S = \{$ 

- (b) Draw the residual network. You can use the below figure (a dotted edge just indicates that in the original graph there is a directed edge at this place).
- (c) Prove that f is a maximum flow by determining a cut  $S \subseteq V(G)$  with cap(S) = val(f) (you might want to recall the proof of the maxflow-mincut theorem for that); provide S by listing its vertices below.

