## Institute for Theoretical Computer Science

## Computational Geometry

URL: http://www.ti.inf.ethz.ch/ew/courses/CG08/

## Exercise 1

You are given

- a star-shaped polygon $P \subset \mathbb{R}^{2}$, represented as a doubly connected list of its vertices $V(P)$,
- and a point $c \in P$ (not necessarily in $V(P)$ ), such that for all $p \in P$ the line segment $\overline{c p}$ is contained in $P$.

Describe an algorithm which triangulates $P$ in linear time. The algorithm could for example output all edges of the triangulation, that are not already edges of the polygon.

## Exercise 2

Let $L$ be a set of $n$ lines in $\mathbb{R}^{2}$ no three of which pass through a common point. Suppose that all lines from $P \subseteq L$ are parallel to each other, no two lines from $L \backslash P$ are parallel to each other, and no line from $L \backslash P$ is parallel to those from $P$. Determine the number of vertices, edges, and faces of the arrangement $\mathcal{A}(\mathrm{L})$ in terms of $n$ and $k:=|\mathrm{P}|$.

## Exercise 3

For an arrangement $\mathcal{A}$ of a set of $n$ lines in $\mathbb{R}^{2}$, let $\mathcal{F}:=\bigcup_{C \text { is cell of } \mathcal{A}} \overline{\mathrm{C}}$ denote the union of the closure of all bounded cells. Show that the complexity (number of vertices and edges of the arrangement lying on the boundary) of $\mathcal{F}$ is $\mathrm{O}(\mathrm{n})$.

## Exercise 4

Given a set of lines in the plane with no three intersecting in a common point, form a graph G whose vertices are the intersections of the lines, with two vertices adjacent if they appear consecutively on one of the lines. Prove that $\chi(G) \leq 3$. ( $\chi$ is the chromatic number of the graph)

