

## Complete $k$ -partite graphs\_\_\_\_\_

$G$  is a **complete  $k$ -partite graph** if there is a partition  $V_1 \cup \dots \cup V_k = V(G)$  of the vertex set, such that  $uv \in E(G)$  iff  $u$  and  $v$  are in *different* parts of the partition. If  $|V_i| = n_i$ , then  $G$  is denoted by  $K_{n_1, \dots, n_k}$ .

The **Turán graph  $T_{n,r}$**  is the complete  $r$ -partite graph on  $n$  vertices whose partite sets differ in size by at most 1. (All partite sets have size  $\lceil n/r \rceil$  or  $\lfloor n/r \rfloor$ .)

**Lemma** Among  $r$ -colorable graphs the Turán graph is the *unique* graph, which has the most number of edges.

*Proof.* Local change.

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## Turán's Theorem\_\_\_\_\_

The **Turán number  $ex(n, H)$**  of a graph  $H$  is the largest integer  $m$  such that there exists an  $H$ -free\* graph on  $n$  vertices with  $m$  edges.

*Example:* Mantel's Theorem states  $ex(n, K_3) = \lfloor \frac{n^2}{4} \rfloor$ .

**Theorem.** (Turán, 1941)

$$ex(n, K_r) = e(T_{n,r-1}) = \left(1 - \frac{1}{r-1}\right) \binom{n}{2} + O(n).$$

*Proof.* Prove by induction on  $r$  that

$G \not\supseteq K_r \implies$  **there is** an  $(r-1)$ -partite graph  $H$  with  $V(H) = V(G)$  and  $e(H) \geq e(G)$ .

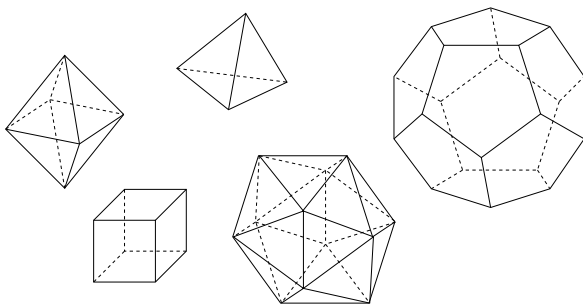
Then apply the Lemma to finish the proof.

\*Here  $H$ -free means that there is no subgraph isomorphic to  $H$

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## Turán-type problems\_\_\_\_\_

**Question.** (Turán, 1941) What happens if instead of  $K_4$ , which is the graph of the tetrahedron, we forbid the graph of some other platonic polyhedra? How many edges can a graph without an octahedron (or cube, or dodecahedron or icosahedron) have?



The platonic solids

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## Erdős-Simonovits-Stone Theorem\_\_\_\_\_

**Theorem.** (Erdős, Simonovits, Stone) For any graph  $H$ ,

$$ex(n, H) = \left(1 - \frac{1}{\chi(H) - 1}\right) \binom{n}{2} + o(n^2).$$

**Corollaries.**

$$\begin{aligned} ex(n, \text{octahedron}) &= \frac{n^2}{4} + o(n^2) \\ ex(n, \text{dodecahedron}) &= \frac{n^2}{4} + o(n^2) \\ ex(n, \text{icosahedron}) &= \frac{n^2}{3} + o(n^2) \\ ex(n, \text{cube}) &= o(n^2) \end{aligned}$$

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## Open problems and Conjectures\_\_\_\_\_

### Known results.

$$\begin{aligned}\Omega(n^{3/2}) &\leq ex(n, Q_3) \leq O(n^{8/5}) \\ \Omega(n^{9/8}) &\leq ex(n, C_8) \leq O(n^{5/4}) \\ \Omega(n^{5/3}) &\leq ex(n, K_{4,4}) \leq O(n^{7/4})\end{aligned}$$

### Conjectures.

$$\begin{aligned}ex(n, K_{t,s}) &= \Theta\left(n^{2-\frac{1}{\min\{t,s\}}}\right) \text{ true for } s = 2, 3 \text{ and } t \geq s \\ &\quad \text{or } s \geq 4 \text{ and } t > (s-1)! \\ ex(n, C_{2k}) &= \Theta\left(n^{1+\frac{1}{k}}\right) \text{ true for } k = 2, 3 \text{ and } 5. \\ ex(n, Q_3) &= \Theta\left(n^{\frac{8}{5}}\right)\end{aligned}$$

If  $H$  is a  $d$ -degenerate bipartite graph, then

$$ex(n, H) = O\left(n^{2-\frac{1}{d}}\right).$$

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## Proof of the ESS Theorem\_\_\_\_\_

**Theorem.** (Erdős-Stone, 1946) For arbitrary fixed integers  $r \geq 2$  and  $t \geq 1$

$$ex(n, T_{rt,r}) = \left(1 - \frac{1}{r-1}\right) \binom{n}{2} + o(n^2).$$

*Proof of ESS from Erdős-Stone Theorem.*  
(Erdős-Simonovits, 1966)

Let  $r = \chi(H)$ .

- $\chi(T_{n,r-1}) < \chi(H)$ , so  $e(T_{n,r-1}) \leq ex(n, H)$ .
- $T_{r\alpha,r} \supseteq H$ , so  $ex(n, T_{r\alpha,r}) \geq ex(n, H)$ , where  $\alpha$  is a constant depending on  $H$ ; say  $\alpha = \alpha(H)$ .

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## Proof of the Erdős-Stone Theorem\_\_\_\_\_

**Erdős-Stone Theorem.** (Reformulation) For any  $\epsilon > 0$  and integers  $r \geq 2, t \geq 1$  there exists an integer  $M = M(r, t, \epsilon)$ , such that any graph  $G$  on  $n \geq M$  vertices with more than  $\left(1 - \frac{1}{r-1} + \epsilon\right) \binom{n}{2}$  edges contains  $T_{rt,r}$ .

We derive this through the following statement.

**Seemingly Weaker Theorem.** For any  $\epsilon > 0$  and integers  $r \geq 2, t \geq 1$  there exists an integer  $N = N(r, t, \epsilon)$ , such that any graph  $G$  on  $n \geq N$  vertices and with  $\delta(G) \geq \left(1 - \frac{1}{r-1} + \epsilon\right) n$  contains  $T_{rt,r}$ .

Note that w.l.o.g.  $\epsilon < \frac{1}{r-1}$ .

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*Derivation of the Erdős-Stone Theorem from the Seemingly Weaker Theorem.*

Let  $G$  be a graph on  $n \geq M(r, t, \epsilon)$ \* vertices with more than  $\left(1 - \frac{1}{r-1} + \epsilon\right) \binom{n}{2}$  edges. Recursively delete vertices which are adjacent to less than  $\left(1 - \frac{1}{r-1} + \frac{\epsilon}{2}\right)$ -fraction of the remaining vertices.

What is the number  $n'$  of vertices we are left with?

We deleted at most  $\sum_{j=n'+1}^n j \left(1 - \frac{1}{r-1} + \frac{\epsilon}{2}\right)$  edges. So

$$e(G) \leq \left(\binom{n+1}{2} - \binom{n'+1}{2}\right) \left(1 - \frac{1}{r-1} + \frac{\epsilon}{2}\right) + \binom{n'}{2}.$$

This implies

$$\frac{\epsilon}{2} \binom{n}{2} - n \leq \left(\frac{1}{r-1} - \frac{\epsilon}{2}\right) \binom{n'}{2} - n'.$$

We choose  $M(r, t, \epsilon)$  such that  $n \geq M(r, t, \epsilon)$  implies  $n' \geq N(r, t, \epsilon/2)$ .

\*At this point we don't know  $M(r, t, \epsilon)$  yet!!! We'll define it in the proof through  $N(r, t, \epsilon/2)$ . (which is known!)

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*Proof of the Seemingly Weaker Theorem.*

Induction on  $r$ .

For  $r = 2$  the claim is true provided  $\frac{\binom{en}{t}}{\binom{n}{t}} > t - 1$ , which is certainly true from some threshold  $N(2, t, \epsilon)$ .

Let  $r \geq 2$  and  $G$  be a graph on  $n \geq N(r + 1, t, \epsilon)^*$  vertices with  $\delta(G) \geq \left(1 - \frac{1}{r} + \epsilon\right)n$ .

We would like to find a  $T_{(r+1)t, r+1}$  in  $G$ .

Let  $s = \left\lceil \frac{t}{\epsilon} \right\rceil$ . By the induction hypothesis<sup>†</sup> there is a  $T_{rs, r}$  in  $G$  with vertex-set  $A_1 \cup \dots \cup A_r$ , where  $|A_1| = \dots = |A_r| = s$ .

$U = V(G) \setminus (A_1 \cup \dots \cup A_r)$ .

$W = \{w \in U : |N(w) \cap A_i| \geq t, i = 1, \dots, r\}$  is the set of vertices eligible to extend some part of  $A_1, \dots, A_r$  into a  $T_{(r+1)t, r+1}$ .

\*Again, we don't know  $N(r + 1, t, \epsilon)$  yet.

†Here we assume  $N(r + 1, t, \epsilon) \geq N(r, s, \epsilon)$ .

Double-count the number of edges missing between  $U$  and  $A_1 \cup \dots \cup A_r$ . They are

- at least  $(|U| - |W|)(s - t)$  and
- at most  $rs \left(\frac{1}{r} - \epsilon\right)n$ .

From this we have

$$|W| \geq \frac{(r-1)\epsilon}{1-\epsilon}n - rs$$

Thus if  $n$  is large enough\* then

$$|W| > \binom{s}{t}^r (t-1).$$

So we can select  $t$  vertices from  $W$ , which are adjacent to the same  $t$  vertices in each  $A_i$ .

\*If  $N(r + 1, t, \epsilon) > \left(\binom{s}{t}\right)^r (t-1) + rs \frac{1-\epsilon}{(r-1)\epsilon}$