

## Bipartite graphs

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A **bipartition** of  $G$  is a specification of two disjoint independent sets in  $G$  whose union is  $V(G)$ .

**Theorem.** (König, 1936) A multigraph  $G$  is bipartite **iff**  $G$  does not contain an odd cycle.

*Proof.*

$\Rightarrow$  Easy.

$\Leftarrow$  Fix a vertex  $v \in V(G)$ . Define sets

$$A := \{w \in V(G) : \exists \text{ an odd } v, w\text{-path}\}$$

$$B := \{w \in V(G) : \exists \text{ an even } v, w\text{-path}\}$$

Prove that  $A$  and  $B$  form a bipartition.

**Lemma.** Every closed odd walk contains an odd cycle.

*Proof.* Strong induction.

## Eulerian circuits

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A multigraph is **Eulerian** if it has a closed trail containing all its edges. A multigraph is called **even** if all of its vertices have even degree.

**Theorem.** Let  $G$  be a connected multigraph. Then

**$G$  is Eulerian iff  $G$  is even.**

*Proof.*

$\Rightarrow$  Easy.

$\Leftarrow$  (Strong) induction on the number of edges.

**Lemma.** If every vertex of a multigraph  $G$  has degree at least 2, then  $G$  contains a cycle.

*Proof.* Extremality: Consider a maximal path...

**Corollary of the proof.** Every even multigraph decomposes into cycles.

## Eulerian trails

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**Theorem.** A connected graph with exactly  $2k$  vertices of odd degree decomposes into  $\max\{k, 1\}$  trails.

*Proof.* Reduce it to the characterization of Eulerian graphs by introducing auxiliary edges.

*Example.* The “little house” can be drawn with one continuous motion.

*Remark.* The theorem is “best possible”, i.e. a decomposition into *less* than  $\max\{k, 1\}$  trails is not possible.

## Proof techniques\_\_\_\_\_

- (Strong) induction
- Extremality
- Double counting

## Neighborhoods and degrees...\_\_\_\_\_

The **neighborhood** of  $v$  in  $G$  is

$$N_G(v) = \{w \in V(G) : vw \in E(G)\}.$$

The **degree** of a vertex  $v$  in graph  $G$  is

$$d_G(v) = |N_G(v)|.$$

The maximum degree of  $G$  is  $\Delta(G) = \max_{v \in V(G)} d(v)$

The minimum degree of  $G$  is  $\delta(G) = \min_{v \in V(G)} d(v)$

$G$  is **regular** if  $\Delta(G) = \delta(G)$

$G$  is  **$k$ -regular** if the degree of each vertex is  $k$ .

The **order** of graph  $G$  is  $n(G) = |V(G)|$ .

The **size** of graph  $G$  is  $e(G) = |E(G)|$ .

## Double counting and bijections I \_\_\_\_\_

**Handshaking Lemma.** For any graph  $G$ ,

$$\sum_{v \in V(G)} d(v) = 2e(G).$$

**Corollary.** Every graph has an **even number** of vertices of **odd degree**.

**No** graph of odd order is regular with odd degree.

**Corollary.** In a graph  $G$  the average degree is  $\frac{2e(G)}{n(G)}$  and hence  $\delta(G) \leq \frac{2e(G)}{n(G)} \leq \Delta(G)$ .

**Corollary.** A  $k$ -regular graph with  $n$  vertices has  $kn/2$  edges.

The  $k$ -dimensional hypercube  $Q_k$ \_\_\_\_\_

$$V(Q_k) = \{0, 1\}^k$$

$$E(Q_k) = \{xy : x \text{ and } y \text{ differ in exactly one coordinate}\}$$

*Properties.*

- $n(Q_k) = 2^k$
- $Q_k$  is  $k$ -regular
- $e(Q_k) = k2^{k-1}$
- $Q_k$  is bipartite
- The number of  $j$ -dimensional subcubes (subgraphs isomorphic to  $Q_j$ ) of  $Q_k$  is  $\binom{k}{j} 2^{k-j}$ .

## Double counting and bijections II\_\_\_\_\_

**Proposition.** Let  $G$  be  $k$ -regular bipartite graph with partite sets  $A$  and  $B$ ,  $k > 0$ . Then  $|A| = |B|$ .

*Proof.* Double count the edges of  $G$ .

**Claim.** The Petersen graph contains ten 6-cycles.

*Proof.* Bijection between 6-cycles and claws. (A claw is a  $K_{1,3}$ .)