

Institut für Theoretische Informatik
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Graph Theory

Problem Set 9

Course Webpage: <http://www.ti.inf.ethz.ch/ew/courses/GT03/>

Due Date: January 7, 2004 at the lecture

Exercise 9.1

(Exercise 4.2.21 in the Textbook)

(!) Let G be a $2k$ -edge-connected graph with at most two vertices of odd degree. Prove that G has a k -edge-connected orientation. (Nash-Williams [1960])

Exercise 9.2

(Exercise 4.2.29 in the Textbook)

Given a graph G , let D be the digraph obtained by replacing each edge with two oppositely-directed edges having the same endpoints (thus D is the symmetric digraph with underlying graph G). Assume that, for all $x, y \in V(D)$, $\kappa_D(x, y) = \lambda_D(x, y)$ holds when $(x, y) \notin E(D)$. Use this hypothesis to prove that also $\kappa_G(x, y) = \lambda_G(x, y)$ for $\{x, y\} \notin E(G)$.

Exercise 9.3

(Exercise not in the Textbook)

In the eighth grade, playing in the little soccer league, our school (*Áldás*) was in a fierce competition with four others: the *Ady*, the *Fillér*, the *Medve*, and the *Törökvészi* elementary schools. During the year every team played every other team 6 times. Despite my brilliant effort as right mid-fielder, in the middle of the season we were standing without a single win, with twelve losses. Our coach was still very enthusiastic and in order to motivate us, got into some complicated argument about how we can still win at least a share of the championship.

The situation was the following. *Ady* had 5 more games with *Fillér*, 2 more games with *Medve*, and 5 with *Törökvészi*. *Fillér* had 3 more games with *Medve*, and 6 with *Törökvészi*. *Medve* had to play 3 more times with *Törökvészi*. *Ady* was at first place with 8 wins, *Törökvészi* was second with 7, *Fillér* had 6, and even the much despised *Medve* was ahead of us with 3 wins.

Was our coach right about his calculations or did he just want to fire us up before our usual showdown with *Medve*? (Hint: you can try to model the problem as a network flow.)

Exercise 9.4

(Exercise not in the Textbook)

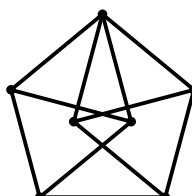
(+) Let $\mathcal{A} = \{A_1, \dots, A_m\}$ and $\mathcal{B} = \{B_1, \dots, B_m\}$ be collections of subsets of a set Y . A **common system of distinct representatives** (CSDR) for \mathcal{A} and \mathcal{B} is a set of distinct elements x_1, \dots, x_m in Y which is both an SDR for \mathcal{A} and an SDR for \mathcal{B} . (For the definition of an SDR, see Exercise 5.5.) Prove that there exists a CSDR for \mathcal{A} and \mathcal{B} if and only if

$$\left| \left(\bigcup_{i \in I} A_i \right) \cap \left(\bigcup_{j \in J} B_j \right) \right| \geq |I| + |J| - m \quad \text{for each pair } I, J \subseteq [m].$$

Exercise 9.5

(Exercise 5.1.1 in the Textbook)

(-) Compute the clique number, the independence number and the chromatic number of the graph below. Is the graph color-critical?



Exercise 9.6

(Exercise 5.1.20 in the Textbook)

(!) Let G be a graph whose odd cycles are pairwise intersecting, meaning that every two odd cycles in G have a common vertex. Prove that $\chi(G) \leq 5$.

Exercise 9.7

(Exercise 5.1.41 in the Textbook)

(!) Prove that $\chi(G) + \chi(\overline{G}) \leq n(G) + 1$. (Nordhaus–Gaddum [1956])

Have a nice Christmas and a happy new year!