## Solution for Assignment 10

- a) Consider the matrix M := AB. We claim that it has rank(M) = n. To see this, observe that rank(B) = n implies C(B) = ℝ<sup>n</sup> because n is also the number of rows of B. Hence, we get C(M) = C(A) (and therefore rank(M) = rank(A) = n). Finally, we can use Proposition 5.5.9 to get (AB)<sup>†</sup> = M<sup>†</sup> = B<sup>†</sup>A<sup>†</sup>.
  - **b)** Let A = CR be the CR decomposition of A with  $C \in \mathbb{R}^{m \times r}$  and  $R \in \mathbb{R}^{r \times n}$  where  $r = \operatorname{rank}(A)$ . Observe that C has full column rank and that R has full row rank. Using the definition of the pseudoinverse, we compute

$$A^{\dagger}AA^{\dagger} = (CR)^{\dagger}CR(CR)^{\dagger} = R^{\dagger}(C^{\dagger}C)(RR^{\dagger})C^{\dagger} = R^{\dagger}C^{\dagger} = A^{\dagger}$$

where we used that  $R^{\dagger}$  is a right inverse of R and  $C^{\dagger}$  a left inverse of C.

c) Assume first that A has full column rank  $n = \operatorname{rank}(A)$ . In this case, we have  $A^{\dagger} = (A^{\top}A)^{-1}A^{\top}$  by definition of the pseudoinverse for matrices with full column rank. Moreover, notice that  $A^{\top}$  has full row rank and hence we also get  $(A^{\top})^{\dagger} = A(A^{\top}A)^{-1}$  by definition of the pseudoinverse for matrices with full row rank. Hence, we get

$$(A^{\dagger})^{\top} = ((A^{\top}A)^{-1}A^{\top})^{\top} = A((A^{\top}A)^{-1})^{\top} = A((A^{\top}A)^{\top})^{-1} = A(A^{\top}A)^{-1} = (A^{\top})^{\dagger}.$$

We conclude that the statements holds for all matrices with full column rank.

Analogously, we can prove that the statement holds if A has full row rank  $m = \operatorname{rank}(A)$ . In that case, we have  $A^{\dagger} = A^{\top} (AA^{\top})^{-1}$  and  $(A^{\top})^{\dagger} = (AA^{\top})^{-1}A$ . Hence, we indeed get

$$(A^{\dagger})^{\top} = (A^{\top}(AA^{\top})^{-1})^{\top} = ((AA^{\top})^{-1})^{\top}A = ((AA^{\top})^{\top})^{-1}A = (AA^{\top})^{-1}A = (A^{\top})^{\dagger}.$$

We conclude that the statement holds for all matrices with full row rank.

It remains to prove the general case, i.e. we do not assume anymore that A has full row rank or full column rank. Then by definition, we have  $A^{\dagger} = R^{\dagger}C^{\dagger}$  where A = CR is a CRdecomposition of A. In particular, we have  $C \in \mathbb{R}^{m \times r}$  and  $R \in \mathbb{R}^{r \times n}$  where  $r = \operatorname{rank}(A)$ . Now observe that we also have  $A^{\top} = R^{\top}C^{\top}$  with  $R^{\top} \in \mathbb{R}^{n \times r}$  and  $C^{\top} \in \mathbb{R}^{r \times m}$  and of course,  $r = \operatorname{rank}(A) = \operatorname{rank}(A^{\top})$ . Hence, we can use Proposition 5.5.9 to get  $(A^{\top})^{\dagger} = (C^{\top})^{\dagger}(R^{\top})^{\dagger}$ . We conclude that

$$(A^{\top})^{\dagger} = (C^{\top})^{\dagger} (R^{\top})^{\dagger} = (C^{\dagger})^{\top} (R^{\dagger})^{\top} = (R^{\dagger} C^{\dagger})^{\top} = (A^{\dagger})^{\top}$$

by using that C has full column rank and R has full row rank and hence  $(C^{\top})^{\dagger} = (C^{\dagger})^{\top}$  and  $(R^{\top})^{\dagger} = (R^{\dagger})^{\top}$ .

**d**) Let A = CR be a CR decomposition of A with  $C \in \mathbb{R}^{m \times r}$  and  $R \in \mathbb{R}^{r \times n}$  where  $r = \operatorname{rank}(A)$ . We can rewrite

$$A^{\dagger}A = (CR)^{\dagger}CR = R^{\dagger}C^{\dagger}CR \stackrel{\text{Prop. 5.5.2}}{=} R^{\dagger}IR = R^{\top}(RR^{\top})^{-1}R$$

and hence we conclude symmetry of  $AA^{\dagger}$  since

$$(A^{\dagger}A)^{\top} = (R^{\top}(RR^{\top})^{-1}R)^{\top} = R^{\top}((RR^{\top})^{-1})^{\top}R = R^{\top}((RR^{\top})^{\top})^{-1}R = R^{\top}(RR^{\top})^{-1}R = A^{\dagger}A.$$

By Theorem 5.2.6, the matrix  $R^{\top}(RR^{\top})^{-1}R = A^{\dagger}A$  is exactly the projection matrix onto the subspace  $\mathbf{C}(R^{\top}) = \mathbf{R}(R) = \mathbf{R}(A) = \mathbf{C}(A^{\top})$  (the equality  $\mathbf{R}(R) = \mathbf{R}(A)$  is due to the observation that R can be obtained from A through row operations and deleting 0-rows, and by recalling that row operations preserve the row space).

- 2. We provide two solutions.
  - In this first solution, we solve this by using our knowledge on pseudoinverses. Consider the function f<sup>-1</sup>: C(A) → C(A<sup>T</sup>) given by f<sup>-1</sup>(x) = A<sup>†</sup>x for all x ∈ C(A). Observe that the composition f<sup>-1</sup> ∘ f is the identity: we know from Exercise 1 that A<sup>†</sup>A is the projection matrix that projects vectors onto the subspace C(A<sup>T</sup>), and hence we have

$$f^{-1}(f(\mathbf{x})) = A^{\dagger}A\mathbf{x} = \mathbf{x}$$

for all  $\mathbf{x} \in \mathbf{C}(A^{\top})$ . This already implies that f is injective. Observe that with an analogous argument we get

$$f(f^{-1}(\mathbf{x})) = AA^{\dagger}\mathbf{x} = \mathbf{x}$$

for all  $\mathbf{x} \in \mathbf{C}(A)$ . Hence,  $f^{-1}$  is injective as well which implies that both f and  $f^{-1}$  are bijective.

Note that the matrix  $A^{\dagger}A$  is in general not the identity matrix. It is crucial that the function f is only defined on  $\mathbf{C}(A^{\top})$  and not on all of  $\mathbb{R}^n$ .

• In this second solution, we start by proving injectivity. For this, let  $\mathbf{x}_1, \mathbf{x}_2 \in \mathbf{C}(A^{\top})$  be arbitrary and assume that  $f(\mathbf{x}_1) = f(\mathbf{x}_2)$ . We want to argue that this implies  $\mathbf{x}_1 = \mathbf{x}_2$ . Observe that we have

$$0 = f(\mathbf{x}_1) - f(\mathbf{x}_2) = A(\mathbf{x}_1 - \mathbf{x}_2)$$

and therefore  $\mathbf{x}_1 - \mathbf{x}_2 \in \mathbf{N}(A)$ . Together with  $\mathbf{C}(A^{\top}) \cap \mathbf{N}(A) = \{0\}$  and  $\mathbf{x}_1 - \mathbf{x}_2 \in \mathbf{C}(A^{\top})$ , we conclude  $\mathbf{x}_1 - \mathbf{x}_2 = 0$  and hence  $\mathbf{x}_1 = \mathbf{x}_2$ .

It remains to prove surjectivity. Let  $\mathbf{y} \in \mathbf{C}(A)$  be arbitrary. By Theorem 5.1.10, we know that there exists  $\mathbf{x} \in \mathbf{C}(A^{\top})$  such that  $\{\mathbf{z} \in \mathbb{R}^d : A\mathbf{z} = \mathbf{y}\} = \mathbf{x} + \mathbf{N}(A)$  (the theorem applies because the set is non-empty since  $\mathbf{y} \in \mathbf{C}(A)$ ). In particular, we have  $f(\mathbf{x}) = A\mathbf{x} = \mathbf{y}$ , as desired.

**3.** For every  $k \in \{1, ..., n-1\}$ , we define  $S_{n-k} = \{1, ..., n-k\}$ .

Our strategy is as follows: We first prove inductively that  $\operatorname{proj}_{S_{n-j}}(P)$  is a polyhedron for all  $1 \le j < n$ . This will then allow us to generalize the proof of Lemma 5.6.4 accordingly.

- For the base case j = 1, observe that  $\text{proj}_{S_{n-j}}(P)$  by Theorem 5.6.3.
- Thus, fix now an arbitrary n > j > 1 and assume that  $\text{proj}_{S'_{n-j}}(P)$  is a polyhedron for j' = j 1 (induction hypothesis).
- Under the above assumption, we want to prove that  $\operatorname{proj}_{S_{n-j}}(P)$  is a polyhedron. Indeed, we know that from the induction hypothesis that  $Q \coloneqq \operatorname{proj}_{S'_{n-j}}(P)$  is a polyhedron. Using Theorem 5.6.3 on Q, we hence conclude that  $\operatorname{proj}_{S_{n-j}}(P)$  is a polyhedron as well.

It remains to prove

$$\operatorname{proj}_{S_{n-i}}(P) = \operatorname{proj}_{S_{n-i}}(\operatorname{proj}_{S_{n-k}}(P)).$$

for all indices  $1 \le k < j < n$ . Let j, k be arbitrary such indices. Observe first that the expression  $\operatorname{proj}_{S_{n-j}}(\operatorname{proj}_{S_{n-k}}(P))$  is now valid because we proved that  $\operatorname{proj}_{S_{n-k}}(P)$  is a polyhedron (and projections are defined on polyhedra).

Consider first an arbitrary  $\mathbf{z} \in \operatorname{proj}_{S_{n-j}}(P)$ . By definition of  $\operatorname{proj}_{S_{n-j}}(P)$ , there exist  $x_{n-j+1}, \ldots, x_n \in \mathbb{R}$  such that the vector

$$\begin{bmatrix} z_1 & \dots & z_{n-j} & x_{n-j+1} & \dots & x_n \end{bmatrix}^\top \in \mathbb{R}^n$$

is in P. By definition of  $\operatorname{proj}_{S_{n-k}}(P)$ , this directly implies that the vector

$$\begin{bmatrix} z_1 & \dots & z_{n-j} & x_{n-j+1} & \dots & x_{n-k} \end{bmatrix}^\top \in \mathbb{R}^{n-k}$$

is in  $\operatorname{proj}_{S_{n-k}}(P)$ . Using the definition of  $\operatorname{proj}_{S_{n-i}}(\operatorname{proj}_{S_{n-k}}(P))$ , we conclude that  $\mathbf{z} \in \operatorname{proj}_{S_{n-i}}(\operatorname{proj}_{S_{n-k}}(P))$ .

For the other direction, consider now an arbitrary vector  $\mathbf{z} \in \text{proj}_{S_{n-j}}(\text{proj}_{S_{n-k}}(P))$ . By definition of  $\text{proj}_{S_{n-j}}(\text{proj}_{S_{n-k}}(P))$ , there exist  $x_{n-j+1}, \ldots, x_{n-k} \in \mathbb{R}$  such that the vector

 $\begin{bmatrix} z_1 & \dots & z_{n-j} & x_{n-j+1} & \dots & x_{n-k} \end{bmatrix}^\top \in \mathbb{R}^{n-k}$ 

is in  $\operatorname{proj}_{S_{n-k}}(P)$ . Now using the definition of  $\operatorname{proj}_{S_{n-k}}(P)$ , there must exist  $x_{n-k+1}, \ldots, x_n \in \mathbb{R}$  such that the vector

 $\begin{bmatrix} z_1 & \dots & z_{n-j} & x_{n-j+1} & \dots & x_n \end{bmatrix}^\top \in \mathbb{R}^n$ 

is in P. We conclude that  $\mathbf{z} \in \operatorname{proj}_{S_{n-i}}(P)$  by the definition of  $\operatorname{proj}_{S_{n-i}}(P)$ .

**4.** In the lecture, it was already proven that, given an arbitrary polyhedron  $P \subseteq \mathbb{R}^n$  for some n, we have  $\operatorname{proj}_{S_{n-i}}(P) \subseteq P^{(i)}$  for all  $i \in [n]$  and that  $P^{(1)} \subseteq \operatorname{proj}_{S_{n-1}}(P)$ . We can use this as base case and proceed by induction over i.

Fix an arbitrary i > 1 and assume as induction hypothesis that we have  $P^{(i-1)} = \operatorname{proj}_{S_{n-i+1}}(P)$  for all polyhedra  $P \subseteq \mathbb{R}^n$  where  $n \ge n - i + 1$ .

In the induction step, we want to prove that we also have  $P^{(i)} = \operatorname{proj}_{S_{n-i}}(P)$  for all polyhedra  $P \subseteq \mathbb{R}^n$  with some  $n \geq i$ . Thus, let P be an arbitrary such polyhedron and consider the polyhedron  $P' = \operatorname{proj}_{S_{n-i+1}}(P) \subseteq \mathbb{R}^{n-i+1}$ . Applying the base case for P' yields  $\operatorname{proj}_{S_{n-i}}(P') = P'^{(1)}$ . Further, we know from Lemma 5.6.4 that  $\operatorname{proj}_{S_{n-i}}(P') = \operatorname{proj}_{S_{n-i+1}}(P)) = \operatorname{proj}_{S_{n-i}}(P)$ . It remains to prove  $P^{(i)} \subseteq P'^{(1)}$ . Using the induction hypothesis, we indeed observe that  $P'^{(1)} = (P^{(i-1)})^{(1)} = P^{(i)}$ , where the equality  $(P^{(i-1)})^{(1)} = P^{(i)}$  follows from Definition 5.6.5.

**5.** Assume that  $P_1 = {\mathbf{x} \in \mathbb{R}^2 : A_1 \mathbf{x} \leq \mathbf{b}_1}$  and  $P_2 = {\mathbf{x} \in \mathbb{R}^2 : A_2 \mathbf{x} \leq \mathbf{b}_2}$  for some  $A_1, A_2 \in \mathbb{Q}^{m \times 2}$  and  $\mathbf{b}_1, \mathbf{b}_2 \in \mathbb{Q}^m$  and natural number m (without loss of generality we can achieve that the two polyhedra have the same number of constraints by just repeating some constraints). Observe that the system

$$\begin{bmatrix} A_1 \\ A_2 \end{bmatrix} \mathbf{x} \le \begin{bmatrix} \mathbf{b}_1 \\ \mathbf{b}_2 \end{bmatrix}$$

has no solution by our assumption  $P_1 \cap P_2 = \emptyset$ . Hence, Farkas lemma implies existence of a vector  $\mathbf{y} \in \mathbb{R}^{2m}$  with  $\mathbf{y} \ge \mathbf{0}, \mathbf{y}^{\top} \begin{bmatrix} A_1 \\ A_2 \end{bmatrix} = \mathbf{0}$ , and  $\mathbf{y}^{\top} \begin{bmatrix} \mathbf{b}_1 \\ \mathbf{b}_2 \end{bmatrix} < 0$ . Let  $\mathbf{y}_1, \mathbf{y}_2 \in \mathbb{R}^m$  be such that  $\mathbf{y} = \begin{bmatrix} \mathbf{y}_1 \\ \mathbf{y}_2 \end{bmatrix}$ . Observe that  $\mathbf{y}^{\top} \begin{bmatrix} A_1 \\ A_2 \end{bmatrix} = \mathbf{0}$  can be rewritten as  $\mathbf{y}_1^{\top} A_1 + \mathbf{y}_2^{\top} A_2 = \mathbf{0}$  and hence  $\mathbf{y}_1^{\top} A_1 = -\mathbf{y}_2^{\top} A_2$ . Similarly, we get  $\mathbf{y}_1^{\top} \mathbf{b}_1 < -\mathbf{y}_2^{\top} \mathbf{b}_2$ .

Now define  $\mathbf{v} := \mathbf{y}_1^\top A_1 \in \mathbb{R}^2$  and  $w := \mathbf{y}_1^\top \mathbf{b}_1 \in \mathbb{R}$ . We claim that  $P_1 \subseteq {\mathbf{x} \in \mathbb{R}^2 : \mathbf{x} \cdot \mathbf{v} \leq w}$ and  $P_2 \subseteq {\mathbf{x} \in \mathbb{R}^2 : \mathbf{x} \cdot \mathbf{v} > w}$ . To prove this, let first  $\mathbf{x} \in P_1$  be arbitrary. Then  $A_1\mathbf{x} \leq \mathbf{b}_1$ . Using  $\mathbf{y}_1 \ge \mathbf{0}$ , we hence get  $\mathbf{y}_1^\top A_1 \mathbf{x} \leq \mathbf{y}_1^\top \mathbf{b}_1$  and thus  $\mathbf{x} \in {\mathbf{x} \in \mathbb{R}^2 : \mathbf{x} \cdot \mathbf{v} \leq w}$ , as desired. Thus, let now  $\mathbf{x} \in P_2$  be arbitrary. Then  $A_2\mathbf{x} \leq \mathbf{b}_2$ . By  $\mathbf{y}_2 \ge \mathbf{0}$  we again get  $\mathbf{y}_2^\top A_2\mathbf{x} \leq \mathbf{y}_2^\top \mathbf{b}_2$ . Using our previous observations, we can rewrite this as  $-\mathbf{y}_1^\top A_1 \mathbf{x} \leq \mathbf{y}_2^\top \mathbf{b}_2 < -\mathbf{y}_1^\top \mathbf{b}_1$ . Multiplying both sides with -1 yields  $\mathbf{x} \in {\mathbf{x} \in \mathbb{R}^2 : \mathbf{x} \cdot \mathbf{v} > w}$ , as desired. This concludes the proof.