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Introduction to Topological Data Analysis
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## Interval Persistence Modules

We consider persistence modules over $\mathbb{R}$ of vector spaces over some field $\mathbb{F}$.
In today's lecture we look at some special persistence modules, called interval modules.

Definition 1. A interval module $\mathbb{I}[\mathrm{b}, \mathrm{d}]$ is an persistence module

$$
\mathrm{V}_{\mathrm{a}}=\left\{\begin{array}{ll}
\mathbb{F} & \text { if } \mathrm{a} \in[\mathrm{~b}, \mathrm{~d}], \\
0 & \text { otherwise. }
\end{array} \quad \text { and } \quad v_{\mathrm{a}, \mathrm{a}^{\prime}}= \begin{cases}\text { id } & \mathrm{b} \leq \mathrm{a} \leq \mathrm{a}^{\prime} \leq \mathrm{d} \\
0 & \text { otherwise } .\end{cases}\right.
$$

Similarly, we can define interval modules on open and clopen intervals, denoted by $\mathbb{I}(b, d), \mathbb{I}(b, d]$, and $\mathbb{I}[b, d)$. We write $\mathbb{I}\langle b, d\rangle$ to include all four of these types.

For an interval module we can easily talk about birth and death as we did in persistent homology. If we have a persistent homology module that is (isomorphic to) an interval module, the birth and death correspond to the boundaries $b, d$ of the interval.

Definition 2. A persistence module $\mathbb{U}$ is called pointwise finite dimensional (p.f.d.) if for all $\mathrm{a} \in \mathbb{R}, \mathrm{U}_{\mathrm{a}}$ has finite dimension.

Note that all p.f.d. persistence modules are also q-tame.

Definition 3. Given two persistence modules $\mathbb{U}, \mathbb{V}$, we define their direct sum $\mathbb{U} \oplus \mathbb{V}$ by $(\mathrm{U} \oplus \mathrm{V})_{\mathrm{a}}=\mathrm{U}_{\mathrm{a}} \oplus \mathrm{V}_{\mathrm{a}}$ and $(u \oplus v)_{\mathrm{a}, \mathrm{a}^{\prime}}=\mathrm{u}_{\mathrm{a}, \mathrm{a}^{\prime}} \oplus v_{\mathrm{a}, \mathrm{a}^{\prime}}$.

Here, the direct sum of maps just means applying the respective maps component-wise.

Proposition 4. If $\mathbb{U}_{1}, \mathbb{U}_{2}$ are $\epsilon$-interleaved, and $\mathbb{V}_{1}, \mathbb{V}_{2}$ are $\delta$-interleaved, then $\mathbb{U}_{1} \oplus \mathbb{V}_{1}$ and $\mathbb{U}_{2} \oplus \mathbb{V}_{2}$ are $\max \{\epsilon, \delta\}$-interleaved.

Proof. W.l.o.g. $\epsilon \geq \delta$, so we need to show that they are $\epsilon$-interleaved. Recall that if two persistence modules are $\delta$-interleaved, they are also $\epsilon$-interleaved. Let $\varphi^{u}, \psi^{u}$ be (series of) functions showing that $\mathbb{U}_{1}, \mathbb{U}_{2}$ are $\epsilon$-interleaved. Similarly, let $\varphi^{v}, \psi^{v}$ be (series of) functions showing that $\mathbb{V}_{1}, \mathbb{V}_{2}$ are $\epsilon$-interleaved. Then, $\varphi^{u} \oplus \varphi^{v}, \psi^{u} \oplus \psi^{v}$ show that $\mathbb{U}_{1} \oplus \mathbb{V}_{1}$ and $\mathbb{U}_{2} \oplus \mathbb{V}_{2}$ are $\epsilon$-interleaved.

If we now have a direct sum of interval modules, we can still nicely talk about birth and death: We just look at each interval module in isolation. The following theorem shows that surprisingly most persistence modules can be expressed as direct sums of interval modules.

Theorem 5 (Structure theorem). Any p.f.d. persistence module decomposes uniquely into interval modules, i.e., we have

$$
\mathbb{U} \cong \bigoplus_{i \in \mathrm{I}} \mathbb{I}\left\langle\mathrm{~b}_{i}, \mathrm{~d}_{\mathfrak{i}}\right\rangle
$$

The intervals $\left\langle\mathrm{b}_{\mathrm{i}}, \mathrm{d}_{\mathrm{i}}\right\rangle$ are exactly the barcodes if $\mathbb{U}$ is a persistent homology module.
Note that unless we have some additional tame-ness condition on $\mathbb{U}$, I is not guaranteed to be finite.

Recall that when we talked about persistent homology, we said that there is some consistent global choice of basis for persistent homology groups. That was a consequence of the structure theorem.

Proposition 6. Consider two interval modules $\mathbb{I}_{1}=\mathbb{I}\left\langle\mathrm{b}_{1}, \mathrm{~d}_{1}\right\rangle$ and $\mathbb{I}_{2}=\mathbb{I}\left\langle\mathrm{b}_{2}, \mathrm{~d}_{2}\right\rangle$. Then, $\mathrm{d}_{\mathrm{I}}\left(\mathbb{I}_{1}, \mathbb{I}_{2}\right)=\mathrm{d}_{\mathrm{b}}\left(\mathrm{DgmI}_{1}, \mathrm{DgmI}_{2}\right)$.

Proof. (This proof has not been shown in the lecture and is not relevant for the exam.) To prove that $d_{I}\left(\mathbb{I}_{1}, \mathbb{I}_{2}\right) \geq d_{b}\left(D g m \mathbb{I}_{1}, D g m \mathbb{I}_{2}\right)$, we show that every upper bound on $d_{I}$ is also an upper bound on $d_{b}$ : assume that we have maps $\varphi, \psi$ showing that the two modules are $\epsilon$-interleaved. Then, consider $\psi_{a+\varepsilon} \circ \varphi_{a}=v_{a, a+2 \epsilon}^{1}$, equality holding because $\varphi, \psi$ certify $\epsilon$-interleaving. Consider $a \in\left\langle b_{1}, d_{1}\right\rangle$.

Case 1: $\quad v_{\mathrm{a}, \mathrm{a}+2 \epsilon}^{1}=0$ for all $\mathrm{a} \in\left\langle\mathrm{b}_{1}, \mathrm{~d}_{1}\right\rangle$. Then, $\mathrm{d}_{1}-\mathrm{b}_{1}<2 \epsilon$, and the (infinity-norm) distance of $\left(b_{1}, d_{1}\right)$ to the diagonal is less than $\epsilon$.

Case 2: $\quad v_{a, a+2 \epsilon}^{1}=i d$ for some $a \in\left\langle b_{1}, d_{1}\right\rangle$. Then, $d_{1}-b_{1} \geq 2 \epsilon$. Furthermore, we have $\varphi_{a}(\mathbb{F})=\mathbb{F}$ for all $a \in\left\langle b_{1}, d_{1}-2 \epsilon\right\rangle$. So, for these $a$, we must also have $a+\epsilon \in\left\langle b_{2}, d_{2}\right\rangle$. This tells us that $\left\langle b_{2}, d_{2}\right\rangle$ must "cover" a large part of $\left\langle b_{1}, d_{1}\right\rangle$, namely we get $b_{2} \leq b_{1}+\epsilon$, and $d_{2} \geq d_{1}-\epsilon$. We can now see that $\left|b_{2}-b_{1}\right| \leq \epsilon$ and $\left|d_{2}-d_{1}\right| \leq \epsilon$ : to violate this, $\left\langle\mathrm{b}_{2}, \mathrm{~d}_{2}\right\rangle$ would have to be a larger interval than $\left\langle\mathrm{b}_{1}, \mathrm{~d}_{1}\right\rangle$ (in particular, it would be longer than $2 \epsilon$ ), and we could thus exchange their roles and get that $b_{1} \leq b_{2}+\epsilon$ and $d_{1} \geq d_{2}-\epsilon$. From this, we get that $d_{\infty}\left(\left(b_{1}, d_{1}\right),\left(b_{2}, d_{2}\right)\right) \leq \epsilon$, and thus get the bound on $d_{b}$.

We now prove the other direction, $d_{I}\left(\mathbb{I}_{1}, \mathbb{I}_{2}\right) \leq d_{b}\left(\mathrm{Dgm}_{1}, D g m \mathbb{I}_{2}\right)$. To see this, we show that from every matching whose longest edge is $\epsilon$, we get an $\epsilon$-interleaving.

Case 1: The two off-diagonal points are matched to the diagonal. Then, we get that $d_{i}-b_{i} \leq 2 \epsilon$ for both of them, and thus for all $\epsilon^{\prime}>\epsilon, \mathbb{I}_{1}$ and $\mathbb{I}_{2}$ are $\epsilon^{\prime}$-interleaved with $\varphi, \psi=0$. Thus, $d_{\mathrm{I}} \leq \epsilon$.

Case 2: The points are matched with each other. Then, $\left|b_{2}-b_{1}\right| \leq \epsilon$ and $\left|d_{2}-d_{1}\right| \leq \epsilon$. Taking $\varphi, \psi=$ id we can see that $\mathbb{I}_{1}$ and $\mathbb{I}_{2}$ are $\epsilon$-interleaved. Thus, $d_{\mathrm{I}} \leq \epsilon$.

Corollary 7. Let $\mathbb{U}, \mathbb{V}$ be p.f.d. persistence modules. Then, $\mathrm{d}_{\mathrm{I}}(\mathbb{U}, \mathbb{V}) \leq \mathrm{d}_{\mathrm{b}}(\mathrm{Dgm} \mathbb{U}, \operatorname{Dgm} \mathbb{V})$.
Proof. We apply the structure theorem to write $\mathbb{U}=\oplus_{i \in \mathrm{I}} \mathbb{I}\left\langle\mathrm{b}_{\mathrm{i}}, \mathrm{d}_{\mathrm{i}}\right\rangle \oplus \oplus_{\mathrm{j} \in \mathrm{J}} 0$ and $\mathbb{V}=$ $\oplus_{j \in J} \mathbb{I}\left\langle\mathrm{~b}_{\mathrm{j}}, \mathrm{d}_{\mathrm{j}}\right\rangle \oplus \oplus_{\mathrm{i} \in \mathrm{I}} \mathrm{O}$. From the Bottleneck matching we get a matching between parts making up $\mathbb{U}$ and $\mathbb{V}$. Since the Bottleneck distance is the maximum length of any edge, we have $d_{b}(D g m \mathbb{U}, \operatorname{Dgm} \mathbb{V}) \geq d_{b}\left(\mathrm{DgmI}_{1}, D g m \mathbb{I}_{2}\right)=d_{I}\left(\mathbb{I}_{1}, \mathbb{I}_{2}\right)$ for every two interval modules that were matched together, where we used Theorem 6. Finally, we use Theorem 4 to get the desired statement.

