

Eidgenössische Technische Hochschule Zürich Swiss Federal Institute of Technology Zurich

Institute of Theoretical Computer Science Patrick Schnider

## Introduction to Topological Data Analysis Scribe Notes 18 FS23

Scribe notes by Simon Weber. Please contact me for corrections. Lecture date: May 4, 2023 Last update: Thursday 4<sup>th</sup> May, 2023, 17:18

## **Interval Persistence Modules**

We consider persistence modules over  $\mathbb{R}$  of vector spaces over some field  $\mathbb{F}$ .

In today's lecture we look at some special persistence modules, called interval modules.

Definition 1. A interval module  $\mathbb{I}[b,d]$  is an persistence module

$$V_{\mathfrak{a}} = egin{cases} \mathbb{F} & \textit{if } \mathfrak{a} \in [\mathfrak{b}, \mathfrak{d}], \ \mathfrak{0} & \textit{otherwise}. \end{cases}$$
 and  $v_{\mathfrak{a},\mathfrak{a}'} = egin{cases} \operatorname{id} & \mathfrak{b} \leq \mathfrak{a} \leq \mathfrak{a}' \leq \mathfrak{d}, \ \mathfrak{0} & \textit{otherwise}. \end{cases}$ 

Similarly, we can define interval modules on open and clopen intervals, denoted by  $\mathbb{I}(b, d)$ ,  $\mathbb{I}(b, d]$ , and  $\mathbb{I}[b, d)$ . We write  $\mathbb{I}(b, d)$  to include all four of these types.

For an interval module we can easily talk about birth and death as we did in persistent homology. If we have a persistent homology module that is (isomorphic to) an interval module, the birth and death correspond to the boundaries b, d of the interval.

**Definition 2.** A persistence module  $\mathbb{U}$  is called pointwise finite dimensional (p.f.d.) if for all  $a \in \mathbb{R}$ ,  $U_a$  has finite dimension.

Note that all p.f.d. persistence modules are also q-tame.

**Definition 3.** Given two persistence modules  $\mathbb{U}, \mathbb{V}$ , we define their direct sum  $\mathbb{U} \oplus \mathbb{V}$  by  $(\mathbb{U} \oplus \mathbb{V})_a = \mathbb{U}_a \oplus \mathbb{V}_a$  and  $(\mathfrak{u} \oplus \mathfrak{v})_{a,a'} = \mathfrak{u}_{a,a'} \oplus \mathfrak{v}_{a,a'}$ .

Here, the direct sum of maps just means applying the respective maps component-wise.

**Proposition 4.** If  $\mathbb{U}_1, \mathbb{U}_2$  are  $\epsilon$ -interleaved, and  $\mathbb{V}_1, \mathbb{V}_2$  are  $\delta$ -interleaved, then  $\mathbb{U}_1 \oplus \mathbb{V}_1$  and  $\mathbb{U}_2 \oplus \mathbb{V}_2$  are max $\{\epsilon, \delta\}$ -interleaved.

*Proof.* W.l.o.g.  $\epsilon \geq \delta$ , so we need to show that they are  $\epsilon$ -interleaved. Recall that if two persistence modules are  $\delta$ -interleaved, they are also  $\epsilon$ -interleaved. Let  $\varphi^{u}, \psi^{u}$  be (series of) functions showing that  $\mathbb{U}_{1}, \mathbb{U}_{2}$  are  $\epsilon$ -interleaved. Similarly, let  $\varphi^{v}, \psi^{v}$  be (series of) functions showing that  $\mathbb{V}_{1}, \mathbb{V}_{2}$  are  $\epsilon$ -interleaved. Then,  $\varphi^{u} \oplus \varphi^{v}, \psi^{u} \oplus \psi^{v}$  show that  $\mathbb{U}_{1} \oplus \mathbb{V}_{1}$  and  $\mathbb{U}_{2} \oplus \mathbb{V}_{2}$  are  $\epsilon$ -interleaved.  $\Box$ 

If we now have a direct sum of interval modules, we can still nicely talk about birth and death: We just look at each interval module in isolation. The following theorem shows that surprisingly most persistence modules can be expressed as direct sums of interval modules.

**Theorem 5** (Structure theorem). Any p.f.d. persistence module decomposes uniquely into interval modules, i.e., we have

$$\mathbb{U} \cong \bigoplus_{i \in I} \mathbb{I} \langle b_i, d_i \rangle.$$

The intervals  $\langle b_i, d_i \rangle$  are exactly the barcodes if  $\mathbb{U}$  is a persistent homology module.

Note that unless we have some additional tame-ness condition on  $\mathbb{U}$ , I is not guaranteed to be finite.

Recall that when we talked about persistent homology, we said that there is some consistent global choice of basis for persistent homology groups. That was a consequence of the structure theorem.

**Proposition 6.** Consider two interval modules  $\mathbb{I}_1 = \mathbb{I}\langle b_1, d_1 \rangle$  and  $\mathbb{I}_2 = \mathbb{I}\langle b_2, d_2 \rangle$ . Then,  $d_I(\mathbb{I}_1, \mathbb{I}_2) = d_b(Dgm\mathbb{I}_1, Dgm\mathbb{I}_2)$ .

*Proof.* (This proof has not been shown in the lecture and is not relevant for the exam.) To prove that  $d_{I}(\mathbb{I}_{1},\mathbb{I}_{2}) \geq d_{b}(Dgm\mathbb{I}_{1},Dgm\mathbb{I}_{2})$ , we show that every upper bound on  $d_{I}$  is also an upper bound on  $d_{b}$ : assume that we have maps  $\varphi, \psi$  showing that the two modules are  $\epsilon$ -interleaved. Then, consider  $\psi_{a+\epsilon} \circ \varphi_{a} = \nu_{a,a+2\epsilon}^{1}$ , equality holding because  $\varphi, \psi$  certify  $\epsilon$ -interleaving. Consider  $a \in \langle b_{1}, d_{1} \rangle$ .

**Case 1:**  $\nu_{a,a+2\varepsilon}^1 = 0$  for all  $a \in \langle b_1, d_1 \rangle$ . Then,  $d_1 - b_1 < 2\varepsilon$ , and the (infinity-norm) distance of  $(b_1, d_1)$  to the diagonal is less than  $\varepsilon$ .

**Case 2:**  $\nu_{a,a+2\epsilon}^1 = \text{id for some } a \in \langle b_1, d_1 \rangle$ . Then,  $d_1 - b_1 \ge 2\epsilon$ . Furthermore, we have  $\varphi_a(\mathbb{F}) = \mathbb{F}$  for all  $a \in \langle b_1, d_1 - 2\epsilon \rangle$ . So, for these a, we must also have  $a + \epsilon \in \langle b_2, d_2 \rangle$ . This tells us that  $\langle b_2, d_2 \rangle$  must "cover" a large part of  $\langle b_1, d_1 \rangle$ , namely we get  $b_2 \le b_1 + \epsilon$ , and  $d_2 \ge d_1 - \epsilon$ . We can now see that  $|b_2 - b_1| \le \epsilon$  and  $|d_2 - d_1| \le \epsilon$ : to violate this,  $\langle b_2, d_2 \rangle$  would have to be a larger interval than  $\langle b_1, d_1 \rangle$  (in particular, it would be longer than  $2\epsilon$ ), and we could thus exchange their roles and get that  $b_1 \le b_2 + \epsilon$  and  $d_1 \ge d_2 - \epsilon$ . From this, we get that  $d_{\infty}((b_1, d_1), (b_2, d_2)) \le \epsilon$ , and thus get the bound on  $d_b$ .

We now prove the other direction,  $d_I(\mathbb{I}_1, \mathbb{I}_2) \leq d_b(Dgm\mathbb{I}_1, Dgm\mathbb{I}_2)$ . To see this, we show that from every matching whose longest edge is  $\epsilon$ , we get an  $\epsilon$ -interleaving.

**Case 1:** The two off-diagonal points are matched to the diagonal. Then, we get that  $d_i - b_i \leq 2\varepsilon$  for both of them, and thus for all  $\varepsilon' > \varepsilon$ ,  $\mathbb{I}_1$  and  $\mathbb{I}_2$  are  $\varepsilon'$ -interleaved with  $\phi, \psi = 0$ . Thus,  $d_I \leq \varepsilon$ .

**Corollary 7.** Let  $\mathbb{U}, \mathbb{V}$  be p.f.d. persistence modules. Then,  $d_{I}(\mathbb{U}, \mathbb{V}) \leq d_{b}(Dgm\mathbb{U}, Dgm\mathbb{V})$ .

*Proof.* We apply the structure theorem to write  $\mathbb{U} = \bigoplus_{i \in I} \mathbb{I}\langle b_i, d_i \rangle \oplus \bigoplus_{j \in J} 0$  and  $\mathbb{V} = \bigoplus_{j \in J} \mathbb{I}\langle b_j, d_j \rangle \oplus \bigoplus_{i \in I} 0$ . From the Bottleneck matching we get a matching between parts making up  $\mathbb{U}$  and  $\mathbb{V}$ . Since the Bottleneck distance is the maximum length of any edge, we have  $d_b(Dgm\mathbb{U}, Dgm\mathbb{V}) \ge d_b(Dgm\mathbb{I}_1, Dgm\mathbb{I}_2) = d_I(\mathbb{I}_1, \mathbb{I}_2)$  for every two interval modules that were matched together, where we used Theorem 6. Finally, we use Theorem 4 to get the desired statement.